

國立成功大學
110學年度碩士班招生考試試題

編 號：190

系 所：電腦與通信工程研究所

科 目：資料結構

日 期：0203

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備 註：不可使用計算機

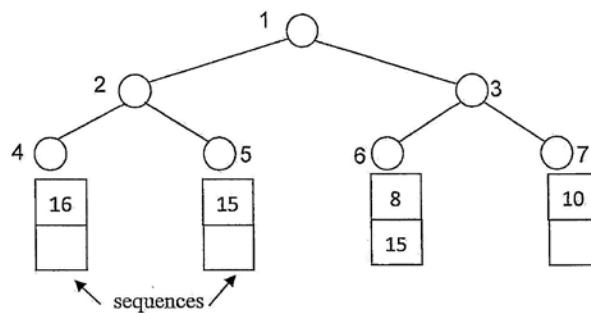
※ 考生請注意：本試題不可使用計算機。請於答案卷(卡)作答，於本試題紙上作答者，不予計分。

1. (30 pts) True or False. Identify all errors in the sentence if it is not correct. You can get the points only when you correct all errors.
- Floyd-Warshall algorithm applies the dynamic programming to find all pair shortest paths and can be applied to a graph with negative weights and negative cycles.
 - There are no back edges, no forward edges, and no cross edges after we apply the breadth-first search algorithm to an undirected graph.
 - A biconnected graph has only one biconnected component and we can use breath-first search algorithm to find biconnected components in a graph.
 - The access time of a hash table with open addressing is related to the load factor.
 - There is no need for parentheses for the prefix notation because there is no ambiguity.
 - There totally exists $n-1$ threads in a post-thread binary tree ^{with} n nodes, where the left thread of a node points to its predecessor and the right thread of the node points to its successor according to the in-order.
 - Dijkstra's algorithm is a greedy algorithm which can be used to find a single source shortest path. It allows a graph $G(V, E)$ with negative weights. The time complexity is $O(E \log V)$ when we use an adjacency matrix to implement it.
 - The failure function of the Knuth-Morris-Pratt algorithm applied to the pattern "abcabcab" is shown as follows:

a	b	c	a	b	c	a	c	a	b
0	-1	-1	0	1	2	0	-1	0	1
 - A circular queue can be used to resolve the problem of necessary to move data if a queue is full, where the queue is implemented by an array with size N . But it can only store $N-1$ data in order to distinguish whether the queue is full or empty. We can pop a data from the location "front" and push a data to the location "rear+1", where "front" and "rear" are pointers to the queue.
 - If $f(n) = \sum_{i=1}^n \log i$, then $f(n) = \theta(\log n)$.

4. (10 pts) Selection tree for minimal values

- a. (2 pts) Show the initial status of the binary tree after the first node in each sequence is processed by a winner tree.
- b. (2 pts) Show the status of the binary tree after the next node is processed by a winner tree.
- c. (2 pts) Show the initial status of the binary tree after the first node in each sequence is processed in a loser tree.
- d. (2 pts) Show the status of the binary tree after the next node is processed by a loser tree.
- e. (2 pts) Compared the two trees, which method is faster? Briefly justify your answer.



5. (18 pts)

- a. (10 pts) Prove the lower bound of the time complexity of the comparison based sorting algorithm for sorting n elements.
- b. (2 pts) Let n denote the number of elements. Let m denote the maximal number of digits for all elements, where the each digit ranges from 0 to k .
What is the time complexity of LSD radix sort algorithm?
- c. (4 pts) Show the result to sort 13, 5, 10, 9, 6 in the non-decreasing order by the LSD radix sort algorithm, where each number is represented by 4 bits binary number. You have to show the results step by step.
- d. (2 pts) Is there exist any conflict on the result in a. and b.? Briefly justify your answer.