## 系所組別：電機資訊學院－資訊聯招

考試科目：程式設計
考試日期：0211，節次：2
第1頁，共3頁
※ 考生請注意：本試題不可使用計算機。 請於答案卷（卡）作答，於本試題紙上作答者，不予計分。 ——Data Structures（50\％）

1．（20\％）
Given a red－black tree in the following figure．


After one node with value 778 is inserted，the resulting red－black tree is shown as follows．

（a）（10\％）Please describe the operation procedures for this insertion．
（b）（ $10 \%$ ）Please indicate the color and value of node $A$ and node $B$ ，respectively．

2．（10\％）Given the expression，$(x+y)^{*} w+u /\left(v+x^{*} w\right)+z$ ，Please show the content in the stack after the operand $v$ is read in postfix transformation．

3．（20\％）
（a）（10\％）Please finish the lost code for the choose function in the following Dijkstra shortest path implementation for a graph without negative－weight edges．

第2頁，共3頁

```
void shortestPath(int v, int cost[][MAX_VERTICES], int distance[], int n, short int found[])
{/*cost is the adjacency matrix*/
    int i,u,w;
    for (i=0; i<n; i++) {
        found[i] = FALSE;
        distance[i] = cost[v][i];
    }
    found [v]= TRUE;
    distance[v]= 0;
    for (i=0; i<n-2; i++){
    u=choose(distance, n, found);
    found[u]= TRUE;
    for (w=0;w<n;w++)
        if (!found[w])
            if (distance[u]+cost[u][w] < distance[w])
                distance[w]=distance[u]+cost[u][w];
    }
}
int choose (int distance[], int n, short int found[])
{
    int i, min, minpos;
    min = INT_MAX;
    minpos=-1;
```

    return minpos;
    \}
（b）（10\％）In a directed graph without a cycle of negative length but with a negative－length edge，we can implement Bellman－Ford algorithm as follows to compute shortest paths．Please fill in the lost code．
void BellmanFord（int $n$ ，int $v$ ）\｛
for（int $\mathrm{i}=0 ; \mathrm{i}<\mathrm{n} ; \mathrm{i}++$ ）
distance $[\mathrm{i}]=\operatorname{cost}[\mathrm{v}][\mathrm{i}]$ ；
for（int $k=2 ; k<=n-1 ; k++$ ）
for（ $\qquad$ ）
for
if（distance［u］＞distance［i］＋cost［i］［u］）
distance［u］＝distance［i］＋cost［i］［u］；

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## 第3頁，共3頁

## 二，Algorithms（50\％）

4．$(10 \%)$ Prove or disprove：The single－source shortest paths problem can be solved in linear time in directed acyclic graphs．
5．（15\％）The matrix－chain multiplication problem can be stated aš follows：Given a chain $\left\langle A_{1}, A_{2}, \ldots, A_{n}\right\rangle$ of $n$ matrices， where for $i=1,2, \ldots, n$ ，matrix $A_{i}$ has dimension $p_{i-1} \times p_{i}$ ，fully parenthesize the product $A_{1} A_{2} \cdots A_{n}$ in a way that minimizes the number of scalar multiplications．Suppose that you have 6 matrices：$A_{1}$ has dimension $30 \times 35, A_{2}$ has dimension $35 \times 15, A_{3}$ has dimension $15 \times 5, A_{4}$ has dimension $5 \times 10, A_{5}$ has dimension $10 \times 20, A_{6}$ has dimension $20 \times 30$ ．Please calculate the minimum number of scalar multiplications．

6．（ $10 \%$ ）Give asymptotic upper and lower bounds for $T(n)=2 T\left(\frac{n}{4}\right)+\sqrt{n}$ ．Assume that $T(n)$ is constant for $n \leq 2$ ． Make your bounds as tight as possible．

7．$(15 \%)$ Consider the problem of finding the 5 －vector $x=\left(x_{i}\right)$ that satisfies

$$
\left(\begin{array}{ccccc}
1 & -1 & 0 & 0 & 0 \\
1 & 0 & 0 & 0 & -1 \\
0 & 1 & 0 & 0 & -1 \\
-1 & 0 & 1 & 0 & 0 \\
-1 & 0 & 0 & 1 & 0 \\
0 & 0 & -1 & 1 & 0 \\
0 & 0 & -1 & 0 & 1 \\
0 & 0 & 0 & -1 & 1
\end{array}\right)\left(\begin{array}{l}
x_{1} \\
x_{2} \\
x_{3} \\
x_{4} \\
x_{5}
\end{array}\right) \leq\left(\begin{array}{c}
0 \\
-1 \\
1 \\
5 \\
-4 \\
-1 \\
-3 \\
-3
\end{array}\right)
$$

Determine whether there exists a solution or there is no solution．

